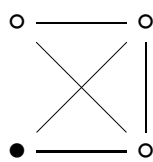
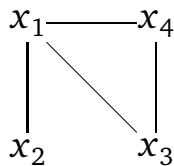


## CLOSURES OF CLANS

The closure properties of clans of the switching classes differ from those of single  $\Delta$ -graphs.

**Example.** The intersection of two clans  $X, Y \in \mathcal{C}(g)$  is always a clan. This is **not so** for switching classes.

Indeed, consider  $\Delta = \mathbb{Z}_2$ , and the next  $g$  together with a switch  $g^\sigma$ :



- $X = \{x_2, x_3, x_4\} \in \mathcal{C}(g)$ ,
- $Y = \{x_1, x_2, x_3\} \in \mathcal{C}(g^\sigma)$ .
- Both in  $\mathcal{C}[g]$ .
- But the intersection  $X \cap Y = \{x_2, x_3\} \notin \mathcal{C}[g]$ .

On the other hand, say that two subsets  $X$  and  $Y$  of a set  $D$  **cross** if they overlap and  $X \cup Y \neq D$ .

**Exercise.** Let  $g$  be a  $\Delta$ -graph with  $X, Y \in \mathcal{C}[g]$  be two crossing clans of  $[g]$ . Show that  $X \cup Y, X \cap Y, X \setminus Y \in \mathcal{C}[g]$ .

## TOWARDS DECOMPOSITION – AGAIN

In analogy to individual  $\Delta$ -graphs, let, for  $[g]$ ,

$$G[g] = (D, E)$$

where

$$E = \{(x, y) \in E_2(D \mid \{x, y\} \in \mathcal{C}[g])\}.$$

**Lemma:**

Let  $\{x, y_1\}, \{x, y_2\}, \{x, y_3\} \in \mathcal{C}[g]$  be different 2-node clans, then  $\{x, y_1, y_2, y_3\}$  induces a complete subgraph of  $G[g]$ .

**Proof. 1.** If  $D = \{x, y_1, y_2, y_3\}$ , then  $\{y_1, y_2\} = D \setminus \{x, y_3\} \in \mathcal{C}[g]$ . Similarly,  $\{y_2, y_3\}, \{y_1, y_3\} \in \mathcal{C}[g]$  and so  $G[g]$  is a complete graph.

**2.** Suppose then  $|D| \geq 5$ .

The clans  $\{x, y_1\}$  and  $\{x, y_2\}$  cross  $\implies \{x, y_1, y_2\} \in \mathcal{C}[g]$ .

The clans  $\{x, y_1, y_2\}$  and  $\{x, y_3\}$  cross  $\implies \{y_1, y_2\} \in \mathcal{C}[g]$ .

Symmetrically,  $\{y_1, y_3\}, \{y_2, y_3\} \in \mathcal{C}[g]$ , and thus the claim. □

## Theorem

Let  $g$  be a  $\Delta^{\delta}$ -graph. If  $G[g]$  is connected, then it is complete or a cycle.

**Proof.** Suppose  $|D| \geq 4$ .

By the lemma, if  $G[g]$  has a vertex of degree  $\geq 3$ , then  $G[g]$  is complete.

Suppose that the degrees are  $\leq 2$ .

By connectivity,  $G[g]$  is either a cycle or a path.

In both cases,  $D = \{x_1, \dots, x_n\}$  so that  $(x_i, x_{i+1})$  are the edges of  $G[g]$ .

Now  $\{x_2, \dots, x_{n-1}\} \in \mathcal{C}[g]$ , since  $\{x_2, \dots, x_k\}$  crosses with  $\{x_k, x_{k+1}\}$  for each  $k = 3, \dots, n-1$ .

Hence  $\{x_1, x_n\} = D \setminus \{x_2, \dots, x_{n-1}\} \in \mathcal{C}[g]$ .

Therefore  $G[g]$  is a (hamiltonian) cycle. □

## QUOTIENTS OF SWITCHING CLASSES

A **factorization** of  $[g]$  is a partition  $D$  into clans  $\mathcal{X} \subseteq \mathcal{C}[g]$ .

These are not quite as straightforward as for the individual  $\Delta$ -graphs, because the clans in  $\mathcal{X}$  may come from **different switches**  $g^\sigma$  of  $g$ .

The **quotient** of  $[g]$  by  $\mathcal{X}$  is

$$[g]/\mathcal{X} = \{h/\mathcal{X} \mid h \in [g], \mathcal{X} \text{ a factorization of } h\}.$$

We shown that  $[g]/\mathcal{X}$  is, **indeed**, a switching class  $[h/\mathcal{X}]$  for some  $h \in [g]$ .

If  $\sigma$  is a constant selector, then  $g \cong g^\sigma$ . Hence

### Lemma

Let  $g$  be a  $\Delta^\delta$ -graph and  $\sigma$  a **constant selector**.

- For each nonempty subset  $X \subseteq D$ , the subgraphs  $g[X]$  and  $g^\sigma[X]$  are isomorphic.
- For each factorization  $\mathcal{X}$  of  $g$ , the quotients  $g/\mathcal{X}$  and  $g^\sigma/\mathcal{X}$ , are isomorphic.

Assume  $\mathcal{X}$  is a factorization of  $g$ .

For a selector  $\sigma: D \rightarrow \Delta$ , let  $\hat{\sigma}: \mathcal{X} \rightarrow D$  be determined by

$$\hat{\sigma}(X) = \sigma(x) \text{ if } x \in X \text{ in } \mathcal{X}.$$

It should be clear that

$$g^\sigma / \mathcal{X} = (g / \mathcal{X})^{\hat{\sigma}}$$

simply because  $g^\sigma / \mathcal{X}$  is isomorphic to the substructure  $g[T]$ , where  $T$  is a transversal of  $\mathcal{X}$ .

Also, the reverse holds: given  $\hat{\sigma}$ , the selector  $\sigma$  can be recovered.

## THEOREM

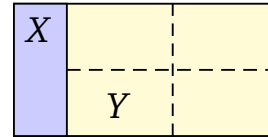
Let  $\mathcal{X}$  be a (proper) factorization of the switching class  $[g]$ . There exists a selector  $\sigma$  such that  $\mathcal{X}$  is a factorization of  $g^\sigma$ , and

$$[g]/\mathcal{X} = [g^\sigma/\mathcal{X}].$$

**Proof.** Denote  $\bar{X} = D \setminus X$  for a  $X \in \mathcal{X}$ .

Isolate  $X$ : there exists  $\sigma$  with  $X, \bar{X} \in \mathcal{C}(g^\sigma)$ .

Let  $Y \in \mathcal{X} \setminus \{X\}$ . Hence  $Y \subseteq \bar{X}$ .



Lemma 7.5(ii): since  $\bar{X} \in \mathcal{C}(g^\sigma)$  and  $Y \in \mathcal{C}((g^\sigma)^\sigma) = \mathcal{C}(g)$  then  $Y \in \mathcal{C}(g^\sigma)$ .

So:  $\mathcal{X}$  is a factorization of  $g^\sigma$  and  $g^\sigma/\mathcal{X}$  is well defined.

Assume that  $\mathcal{X}$  is a factorization of some  $h \in [g]$  (among with  $g^\sigma$ ).

Since  $[g] = [g^\sigma]$ , there is a  $\tau$  with  $h = (g^\sigma)^\tau$ .

- Now, each  $X \in \mathcal{X}$  is a clan in both  $g^\sigma$  and  $(g^\sigma)^\tau$ .  
By Lemma 7.5(i),  $\tau$  is **constant on each**  $X \in \mathcal{X}$ .
- Hence  $h/\mathcal{X} = (g^\sigma/\mathcal{X})^{\hat{\tau}}$ . **Consequently**,  $h/\mathcal{X} \in [g^\sigma/\mathcal{X}]$ .
- **Finally**, let  $\hat{\tau}: \mathcal{X} \rightarrow \Delta$  be given.

By Lemma 7.5(ii),  $\mathcal{X}$  is a factorization of  $g^{\tau\sigma}$ .

Since  $g^{\tau\sigma}/\mathcal{X} = (g^\sigma/\mathcal{X})^{\hat{\tau}}$ ,

$$[g]/\mathcal{X} = \{(g^\sigma/\mathcal{X})^{\hat{\tau}} \mid \hat{\tau} \text{ a selector}\}$$

as required.

□